Chapter 2: Lists

From Theory to Python

List performance

Table from the book Chapter 2.6: Lists

(http://interactivepython.org/runestone/static/pythonds/AlgorithmAnalysis/Lists.html)

Operation	Big-O Efficiency	contains (in)	O(n)
index []	O(1)	get slice [x:y]	O(k)
index assignment	O(1)	del slice	O(n)
append	O(1)	set slice	O(n+k)
pop()	O(1)	reverse	O(n)
pop(i)	O(n)	concatenate	O(k)
insert(i,item)	O(n)	sort	O(n log n)
del operator	O(n)	multiply	O(nk)
iteration	O(n)		O(IIK)

Fast or not?

x = ["a", "b", "c"]
x[2] x[2] = "d" x.append("d") x.insert(0, "d") x[3:5] x.s
ort()

What about len(x) ? If you don't know the answer, try googling it!

Sublist iteration performance

get slice time complexity is O(k), but what about memory? It's the same!

So if you want to iterate a part of a list, beware of slicing! For example, slicing a list like this can occupy much more memory than necessary:

In [2]:

x = range(1000)

print [2*y for y in x[100:200]]

[200, 202, 204, 206, 208, 210, 212, 214, 216, 218, 220, 222, 224, 226, 228, 230, 2 32, 234, 236, 238, 240, 242, 244, 246, 248, 250, 252, 254, 256, 258, 260, 262, 264 , 266, 268, 270, 272, 274, 276, 278, 280, 282, 284, 286, 288, 290, 292, 294, 296, 298, 300, 302, 304, 306, 308, 310, 312, 314, 316, 318, 320, 322, 324, 326, 328, 33 0, 332, 334, 336, 338, 340, 342, 344, 346, 348, 350, 352, 354, 356, 358, 360, 362, 364, 366, 368, 370, 372, 374, 376, 378, 380, 382, 384, 386, 388, 390, 392, 394, 3 96, 398] The reason is that, depending on the Python interpreter you have, slicing like x[100:200] at loop start can create a *new* list. If we want to explicitly tell Python we just want to iterate through the list, we can use the so called <u>itertools (https://docs.python.org/2/library/itertools.html)</u>. In particular, the <u>islice</u> (<u>https://docs.python.org/2/library/itertools.html#itertools.islice</u>)</u> method is handy, with it we can rewrite the list comprehension above like this:

In [3]:

import itertools

```
print [2*y for y in itertools.islice(x, 100, 200)]
```

[200, 202, 204, 206, 208, 210, 212, 214, 216, 218, 220, 222, 224, 226, 228, 230, 2 32, 234, 236, 238, 240, 242, 244, 246, 248, 250, 252, 254, 256, 258, 260, 262, 264 , 266, 268, 270, 272, 274, 276, 278, 280, 282, 284, 286, 288, 290, 292, 294, 296, 298, 300, 302, 304, 306, 308, 310, 312, 314, 316, 318, 320, 322, 324, 326, 328, 33 0, 332, 334, 336, 338, 340, 342, 344, 346, 348, 350, 352, 354, 356, 358, 360, 362, 364, 366, 368, 370, 372, 374, 376, 378, 380, 382, 384, 386, 388, 390, 392, 394, 3 96, 398]

Exercises

Implement swap

Try to code and test the swap function from <u>selection sort (slide 29 theory)</u> (<u>http://disi.unitn.it/~montreso/sp/slides/03-analisi.pdf</u>):

swap(ITEM[]
$$A$$
, int x , int y)
int $temp = A[x]$
 $A[x] = A[y]$
 $A[y] = temp$

• Use the following skeleton to code it

• Check carefully all the test cases, in particular test_swap_property and test_double_swap. They show two important properties of the swap function. Make sure you understand why these tests should succeed.

```
In [4]:
```

```
import unittest
def swap(A, x, y):
    In the array A, swaps the elements at position x and y.
    .....
    raise Exception("TODO implement me!")
class SwapTest(unittest.TestCase):
    def test_one_element(self):
        v = ['a'];
        swap(v,0,0)
        self.assertEqual(v, ['a'])
    def test two elements(self):
        v = ['a', 'b'];
        swap(v,0,1)
        self.assertEqual(v, ['b','a'])
    def test return none(self):
        v = ['a', 'b', 'c', 'd'];
        self.assertEquals(None, swap(v,1,3))
    def test_long_list(self):
        v = ['a', b', c', d'];
        swap(v,1,3)
        self.assertEqual(v, ['a', 'd', 'c', 'b'])
    def test_swap_property(self):
        v = ['a', b', c', d'];
w = ['a', b', c', d'];
        swap(v,1,3)
        swap(w,3,1)
        self.assertEqual(v, w)
    def test double swap(self):
        v = ['a', 'b', 'c', 'd'];
swap(v,1,3)
        swap(v,1,3)
        self.assertEqual(v, ['a','b','c','d'])
```

Implement partial_min_pos

Try to code and test the partial min pos function from <u>selection sort (slide 29 theory)</u> (<u>http://disi.unitn.it/~montreso/sp/slides/03-analisi.pdf</u>):

```
int min(ITEM[] A, int i, int n)

int min = i

for j = i + 1 to n - 1 do

if A[j] < A[min] then

\lim_{k \to 0} min = j

return min
```

% Partial minimum position

% New partial minimum

- Use the following skeleton to code it
- add some test to the provided testcase class

Notice that

- we renamed min to partial_min_pos to avoid name collision with Python standard library min function
- it is not necessary to pass list length n, as it is already stored in Python implementation of lists, and we can retrieve it in 0(1) time with len(A)

In [5]:

import unittest

```
def partial_min_pos(A, i):
```

Return the index of the element in list A which is lesser or equal than all other values i n A

that start from index i included

raise Exception("TODO implement me!")

class PartialMinPosTest(unittest.TestCase):

```
def test_one_element(self):
    self.assertEqual(partial_min_pos([1],0),0)
```

```
def test_two_elements(self):
    self.assertEqual(partial_min_pos([1,2],0),0)
    self.assertEqual(partial_min_pos([2,1],0),1)
    self.assertEqual(partial_min_pos([2,1],1),1)

def test long list(self):
```

```
self.assertEqual(partial_min_pos([8,9,6,5,7],2),3)
```

Implement selection_sort

Try to code and test the selectionSort from <u>selection sort (slide 29 theory)</u> (<u>http://disi.unitn.it/~montreso/sp/slides/03-analisi.pdf</u>):

selectionSort(ITEM[] A , int n)
for $i = 0$ to $n - 2$ do
int $min = min(A, i, n)$
$swap(A, i, \min)$

Use the following skeleton to code it and add some test to the provided testcase class.

Notice that

- we renamed selectionSort to selection_sort because it is a more pythonic name (https://www.python.org/dev/peps/pep-0008/#function-names)
- it is not necessary to pass list length n, as it is already stored in Python implementation of lists, and we can retrieve it in 0(1) time with len(A)
- On the book website, there is <u>an implementation of the selection sort</u> (<u>http://interactivepython.org/runestone/static/pythonds/SortSearch/TheSelectionSort.html</u>) with a nice animated histogram showing a sorting process. Differently from the slides, instead of selecting the minimal element the algorithm on the book selects the *maximal* element and puts it to the right of the array.

In [6]:

```
import unittest
```

```
def selection sort(A):
    Sorts the list A in-place in O(n^2) time.
    raise Exception("TODO implement me!")
class SelectionSortTest(unittest.TestCase):
    def test zero elements(self):
        v = []
        selection sort(v)
        self.assertEqual(v,[])
    def test return none(self):
        self.assertEquals(None, selection sort([2]))
    def test one element(self):
        v = ["a"]
        selection sort(v)
        self.assertEqual(v,["a"])
    def test two elements(self):
        v = [2,1]
        selection sort(v)
        self.assertEqual(v,[1,2])
    def test three elements(self):
        v = [2, 1, 3]
        selection sort(v)
        self.assertEqual(v,[1,2,3])
    def test piccinno list(self):
        v = [23, 34, 55, 32, 7777, 98, 3, 2, 1]
        selection sort(v)
        vcopy = v[:]
        vcopy.sort()
        self.assertEqual(v, vcopy)
```

Implement gap_rec

Try to code and test the gap function from <u>recursion theory slides (slide 21)</u> (<u>http://disi.unitn.it/~montreso/sp/slides/02-recursion.pdf</u>):

ap(int[] L, int i, int j)	
j == i + 1 then	
$_$ return j	
$a = \lfloor (i+j)/2 \rfloor$	
$L[m] \leq L[i]$ then	
return $gap(L, m, j)$	
se	
_ return gap (L, i, m)	

Use the following skeleton to code it and add some test to the provided testcase class. To understand what's going on, try copy pasting your solution in <u>Python tutor (http://pythontutor.com/visualize.html#mode=edit)</u> and hit Visualize execution and then Forward to step through the process

Notice that

- We created a function gap_rec to differentiate it from the iterative one
- Users of gap_rec function might want to call it by passing just a list, in order to find any gap in the whole list. So for convenience the new function gap_rec(L) only accepts a list, without indexes i and j. This function just calls the other function gap_rec_helper that will actually contain the recursive calls. So your task is to translate the pseudocode of gap into the Python code of gap_rec_helper, which takes as input the array and the indexes as gap does. Adding a helper function is a frequent pattern you can find when programming recursive functions.

WARNING: The specification of gap_rec assumes the input is always a list of at least two elements, and that the first element is less or equal than the last one. If these conditions are not met, function behaviour could be completely erroneus!

When preconditions are not met, execution could stop because of an error like index out of bounds, or, even worse, we might get back some wrong index as a gap! To prevent misuse of the function, a good idea can be putting a check at the beginning of the gap_rec function. Such check should immediately stop the execution and raise an error if the parameters don't satisfy the preconditions. One way to do this could be to write some assertion (testing#Assertions) like this:

assert len(L) >= 2
 assert L[0] <= L[len(L)-1]</pre>

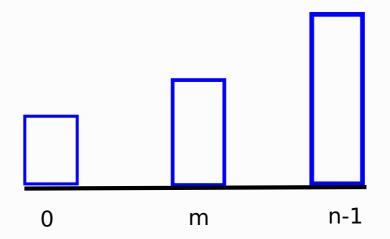
- These commands will make python interrupt execution and throw an error as soon it detects list L is too small or with wrong values
- This kind of behaviour is also called *fail fast*, which is better than returning wrong values!
- You can put any condition you want after assert, but ideally they should be fast to execute.

In [8]:

GOOD PRACTICE: Notice I wrote as a comment what the helper function is expected to receive. Writing down specs often helps understanding what the function is supposed to do, and helps users of your code as well!

COMMANDMENT: You shall also write on paper!

To get an idea of how gap_rec is working, draw histograms on paper like the following, with different heights at index m:



Notice how at each recursive call, we end up with a histogram that is similar to the initial one, that is, it respects the same preconditions (a list of size >= 2 where first element is smaller or equal than the last one)

• Look at the iterative gap here:

```
def gap_iter(L):
    for i in range(1,len(L)):
        if L[i-1] < L[i]:
            return i
        return -1</pre>
```

- What is the complexity of gap_rec? Is it faster or slower than gap_iter ?
- Assuming L contains n >= 2 integers such that L[0] < L[n-1], will the recursive gap always give the same result as the iterative one? If we just change function names, can we run the same test case against both implementations? (Careful!)

```
In [10]:
import unittest
def gap rec(L):
    """ Searches a gap in list L
   Given a list L containing n \ge 2 integers such that L[0] < L[n-1], returns a gap in the li
st.
    A gap is an index i, 0 < i < n such that L[i-1] < L[i]
    return gap_helper(L, 0, len(L)-1)
def gap_helper(L, i, j):
    """ Searches a gap in sublist L[i:j]
    Given a list L containing n \ge 2 integers such that L[i] < L[j], returns a gap in the subl
ist L[i:j]
   A gap is an index z, i < z < j+1 such that L[z-1] < L[z]
    raise Exception("TODO implement me!")
class GapRecTest(unittest.TestCase):
    def test two elements(self):
        self.assertEqual(gap_rec([1,2]),1)
    def test three elements middle(self):
        self.assertEquals(gap_rec([1,3,3]), 1)
    def test three elements right(self):
        self.assertEquals(gap_rec([1,1,3]), 2)
```

Implement binary_search_rec

Try to code and test the binarySearch recursive function from <u>recursion theory slides (slide 21)</u> (<u>http://disi.unitn.it/~montreso/sp/slides/02-recursion.pdf</u>):

$$\begin{split} & \overbrace{\mathbf{int \ binarySearch(\mathbf{int}[] \ L, \ \mathbf{int} \ v, \ \mathbf{int} \ i, \ \mathbf{int} \ j)} \\ & \overbrace{\mathbf{if} \ i > j \ \mathbf{then}} \\ & _ \mathbf{return \ -1} \\ & \mathbf{else} \\ & \overbrace{\mathbf{int} \ m = \lfloor (i+j)/2 \rfloor} \\ & \overbrace{\mathbf{if} \ L[m] = v \ \mathbf{then}} \\ & _ \mathbf{return \ m} \\ & \mathbf{else} \ \mathbf{if} \ L[m] < v \ \mathbf{then} \\ & _ \mathbf{return \ binarySearch(L, v, m+1, j)} \\ & \underbrace{\mathbf{else}} \\ & _ \mathbf{return \ binarySearch(L, v, i, m-1)} \end{split}$$

- Use the following skeleton to code it
- add some test to the provided testcase class
- Does the pseudocode algorithm work with the empty list?
- What happens if we allow non-distinct numbers? Does it work anyway?
- What is the time complexity of the recursive version?
- What is the memory complexity of the recursive version?

Notice that

- we renamed binarySearch to binary_search_rec to have more pythonic name and differentiate it from the iterative one
- the renamed function binary_search_rec(L) only accepts a list, without indexes i and j, we will need a
 way to specify them when we translate the pseudocode. You can follow the same pattern used for
 gap_rec_helper
 - SUGGESTION : write as a comment what the helper function is expected to receive. Can it receive an empty list? Can it receive indices out of bounds? You decide the assumptions, but once they are decided you should make sure that unacceptable values don't get into it!
- To understand what's going on, try copy pasting your solution in <u>Python tutor</u> (<u>http://pythontutor.com/visualize.html#mode=edit</u>) and hit Visualize execution and then Forward to step through the process
- Remember that even experienced programmers tend to fail implementing the binary search at first time, it's easy to get wrong indexes! So good tests here can really spot issues.

In [11]:

```
import unittest
def binary search rec(L,v ):
    """ Searches value v in sorted list L
    Given a list L containing n distinct sorted integers, returns the index position
    of element with value v, or -1 if not found
    raise Exception("TODO implement me!")
class BinarySearchRecTest(unittest.TestCase):
    def test empty(self):
        self.assertEqual(binary search rec([], 7), -1)
    def test one element found(self):
        self.assertEqual(binary search rec([7],7),0)
    def test one element not found(self):
        self.assertEqual(binary search rec([6],7),-1)
    def test one negative element not found(self):
        self.assertEqual(binary search rec([-7],7),-1)
    def test two elements found right(self):
        self.assertEquals(binary_search_rec([6,7],7), 1)
    def test two elements not found(self):
        self.assertEquals(binary search rec([6,7],3), -1)
    def test two elements found left(self):
        self.assertEquals(binary search rec([6,7],6), 0)
    def test long list(self):
        self.assertEquals(binary search rec([2,4,5,7,9],7), 3)
```

Implement binary_search_iter

Try to code and test the iterativeBinarySearch function from <u>Introduction slides (slide 18)</u> (<u>http://disi.unitn.it/~montreso/sp/slides/01-introduzione.pdf</u>):

int iterativeBinarySearch(int[] A, int n, int v)

```
\begin{array}{ll} \operatorname{int} i = 0 & \operatorname{if} i > j \text{ or } A[m] \neq v \text{ then} \\ & | \operatorname{return} -1 \\ & | \operatorname{return} -1 \\ & | \operatorname{return} m \\ \end{array}
\begin{array}{ll} \operatorname{if} A[m] < v \text{ then} \\ & | i = m + 1 \\ & | else \\ & | j = m - 1 \\ & m = \lfloor (i+j)/2 \rfloor \end{array}
```

- This time, there's no code skeleton and you're on your own!
- Try to reuse test cases from the recursive version
- What is the time complexity of the iterative version? Is it different from the recursive version?
- What is the memory complexity of the iterative version? Is it different from the recursive version?

Solutions

swap solution

```
In [12]:
import unittest
def swap(A, x, y):
    In the array A, swaps the elements at position x and y.
    .....
    temp = A[x]
    A[x] = A[y]
    A[y] = temp
class SwapTest(unittest.TestCase):
    def test one element(self):
        v = ['a'];
        swap(v,0,0)
        self.assertEqual(v, ['a'])
    def test two elements(self):
        v = ['a','b'];
        swap(v,0,1)
        self.assertEqual(v, ['b','a'])
    def test return none(self):
        v = ['a','b', 'c', 'd'];
        self.assertEquals(None, swap(v,1,3))
    def test_long_list(self):
        v = ['a', b', c', d'];
        swap(v,1,3)
        self.assertEqual(v, ['a', 'd','c', 'b'])
    def test swap property(self):
        v = ['a', 'b', 'c', 'd'];
w = ['a', 'b', 'c', 'd'];
        swap(v, 1, 3)
        swap(w,3,1)
        self.assertEqual(v, w)
    def test_double_swap(self):
        v = ['a','b','c','d'];
        swap(v,1,3)
        swap(v,1,3)
        self.assertEqual(v, ['a','b','c','d'])
```

partial_min_pos solution

```
In [14]:
```

```
import unittest
def partial min pos(A, i):
    Return the index of the element in list A which is lesser or equal than all other values i
nΑ
    that start from index i included
    .....
    pm = i
    for j in range(i+1, len(A)):
        if (A[j] < A[pm]):
            pm = j
    return pm
class PartialMinPosTest(unittest.TestCase):
    def test one element(self):
        self.assertEqual(partial min pos([1],0),0)
    def test two elements(self):
        self.assertEqual(partial min pos([1,2],0),0)
        self.assertEqual(partial min pos([2,1],0),1)
        self.assertEqual(partial_min_pos([2,1],1),1)
   def test long list(self):
        self.assertEqual(partial_min_pos([8,9,6,5,7],2),3)
```

selection_sort solution

```
In [16]:
```

```
import unittest
def selection sort(A):
    .....
    Sorts the list A in-place in O(n^2) time.
    for i in range(0, len(A)-1):
        m = partial_min_pos(A, i)
        swap(A, i, m)
class SelectionSortTest(unittest.TestCase):
    def test zero elements(self):
        v = []
        selection sort(v)
        self.assertEqual(v,[])
    def test return none(self):
        self.assertEquals(None, selection sort([2]))
    def test one element(self):
        v = ["a"]
        selection sort(v)
        self.assertEqual(v,["a"])
    def test two elements(self):
        v = [2,1]
        selection_sort(v)
        self.assertEqual(v,[1,2])
    def test_three_elements(self):
        v = [2, 1, 3]
        selection_sort(v)
        self.assertEqual(v,[1,2,3])
    def test piccinno list(self):
        v = \overline{[23, 34, 55, 32, 7777, 98, 3, 2, 1]}
        selection sort(v)
        vcopy = v[:]
        vcopy.sort()
        self.assertEqual(v, vcopy)
```

gap_rec solution

```
In [18]:
```

```
import unittest
def gap_rec(L):
    Given a list L containing n \ge 2 integers such that L[0] < L[n-1], returns a gap in the li
st.
    A gap is an index i, 0 < i < n such that L[i-1] < L[i]
    return gap helper(L, 0, len(L)-1)
def gap_helper(L, i, j):
    Given a list L containing n \ge 2 integers such that L[i] < L[j], returns a gap in the subl
ist L[i:j]
    A gap is an index z, i < z < j+1 such that L[z-1] < L[z]
    0.01
    if j == i + 1:
       return j
    m = (i+i) // 2
                     # remember in every python version // operator behaves the same and floor
s the result
    if (L[m] <= L[i]):
        return gap helper(L, m, j)
    else:
        return gap_helper(L, i, m)
class GapRecTest(unittest.TestCase):
    def test_two_elements(self):
        self.assertEqual(gap rec([1,2]),1)
    def test three elements middle(self):
        self.assertEquals(gap_rec([1,3,3]), 1)
    def test three elements right(self):
        self.assertEquals(gap rec([1,1,3]), 2)
```

binary_search_rec solution

In [20]:

import unittest

```
def binary search rec(L,v ):
    """ Searches value v in sorted list L
    Given a list L containing n distinct sorted integers, returns the index position
    of element with value v, or -1 if not found
    .....
    return binary search helper(L,v, 0, len(L)-1)
def binary search helper(L, v, i, j):
    """ Helper for the recursive binary search
    Given a list L containing n distinct sorted integers, returns the index position
    of element with value v if it is present in sublist L[i:j], or -1 if not found
    ......
    if i > j:
        return -1
    m = (i+j) // 2 # remember in every python version // operator behaves the same and floor
s the result
    # print "L = ", L
   # print "v = "
                  , v
   # print "m = ", m
   # print "i = ", i
```

```
# print "j = ", j
    if L[m] == v:
        return m
    elif L[m] < v:</pre>
        return binary search helper(L, v, m + 1, j)
    else:
        return binary search helper(L, v, i, m - 1)
class BinarySearchRecTest(unittest.TestCase):
    def test empty(self):
        self.assertEqual(binary search rec([], 7), -1)
    def test one element found(self):
        self.assertEqual(binary search rec([7],7),0)
    def test one element not found(self):
        self.assertEqual(binary search rec([6],7),-1)
    def test one negative element not found(self):
        self.assertEqual(binary search rec([-7],7),-1)
    def test two elements found right(self):
        self.assertEquals(binary_search_rec([6,7],7), 1)
    def test_two_elements_not_found(self):
        self.assertEquals(binary search rec([6,7],3), -1)
    def test two elements found left(self):
        self.assertEquals(binary search rec([6,7],6), 0)
    def test long list(self):
        self.assertEquals(binary search rec([2,4,5,7,9],7), 3)
    def test not distinct found(self):
        self.assertEquals(binary search rec([7,7],7), 0)
```

```
def test_not_distinct_not_found(self):
    self.assertEquals(binary_search_rec([7,7],5), -1)
```